

(Naive) A^1 - homotopy equivalences

and

Theorems of Whitehead and Zariski

Notation: $X, Y, Z / k$ for varieties over an arbitrary field k
(often X, Y, Z will be smooth, sometimes proper or projective)

Goals: approach A^1 -homotopy theory from a very low-tech point of view and derive some basic results in this setting

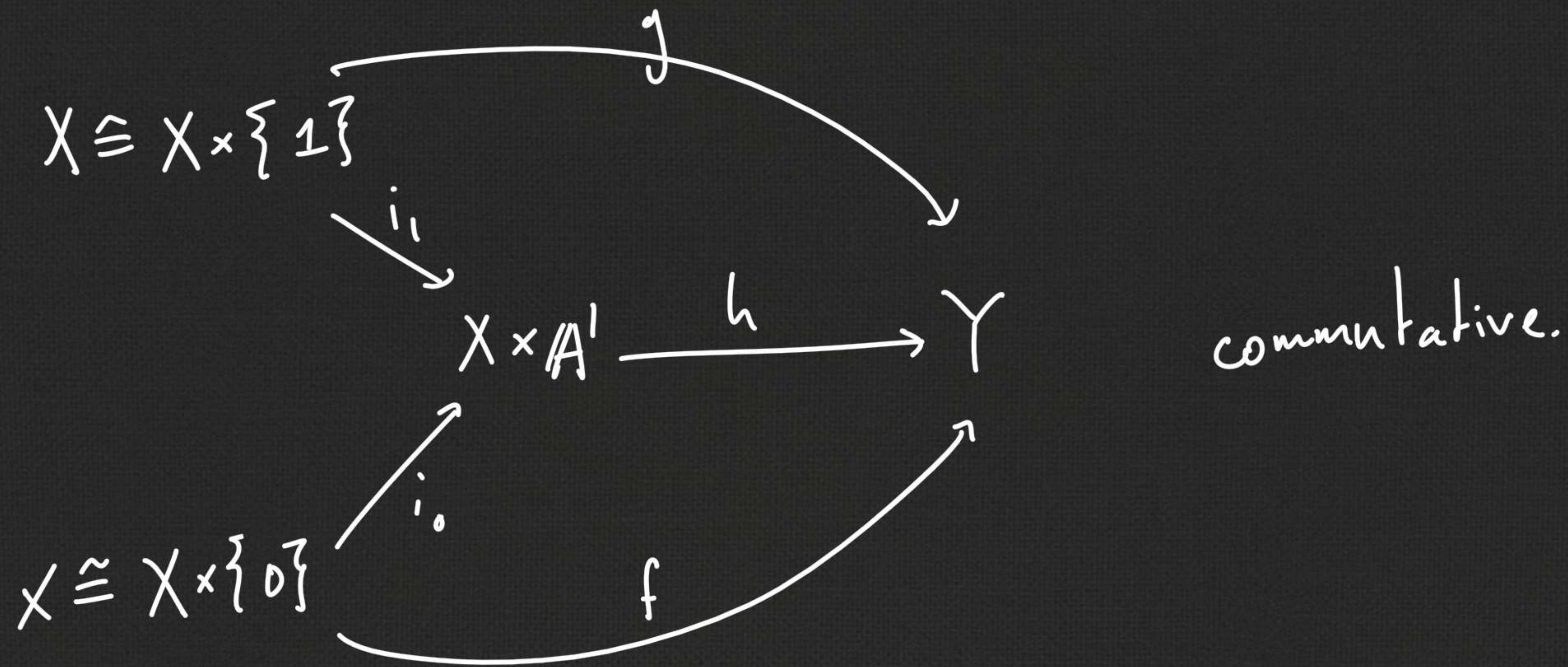
(Feel free to ask questions at any time)

Part 1 :

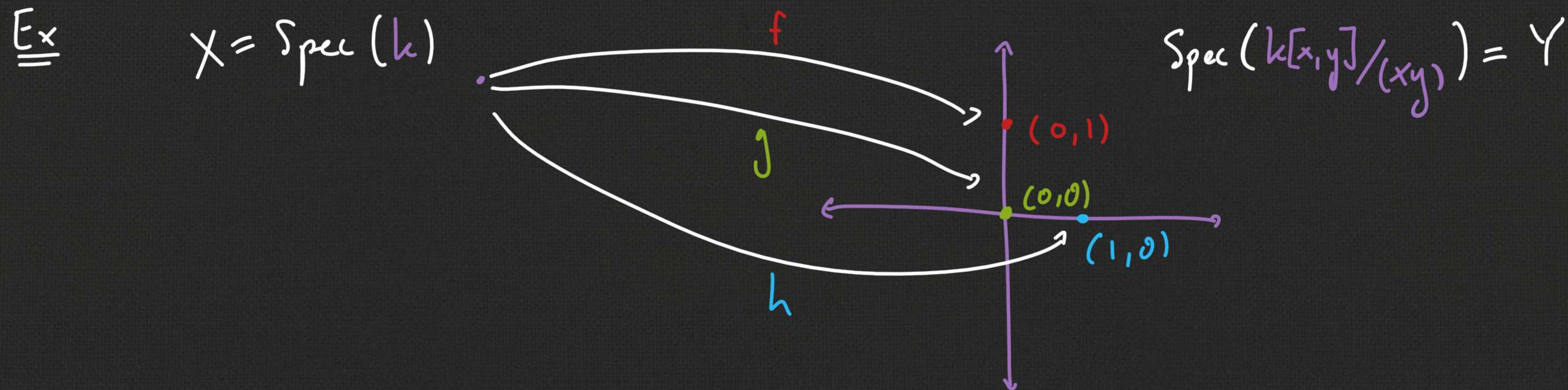
naive \mathbb{A}^1 -homotopy theory

Starting from somewhere simple ...

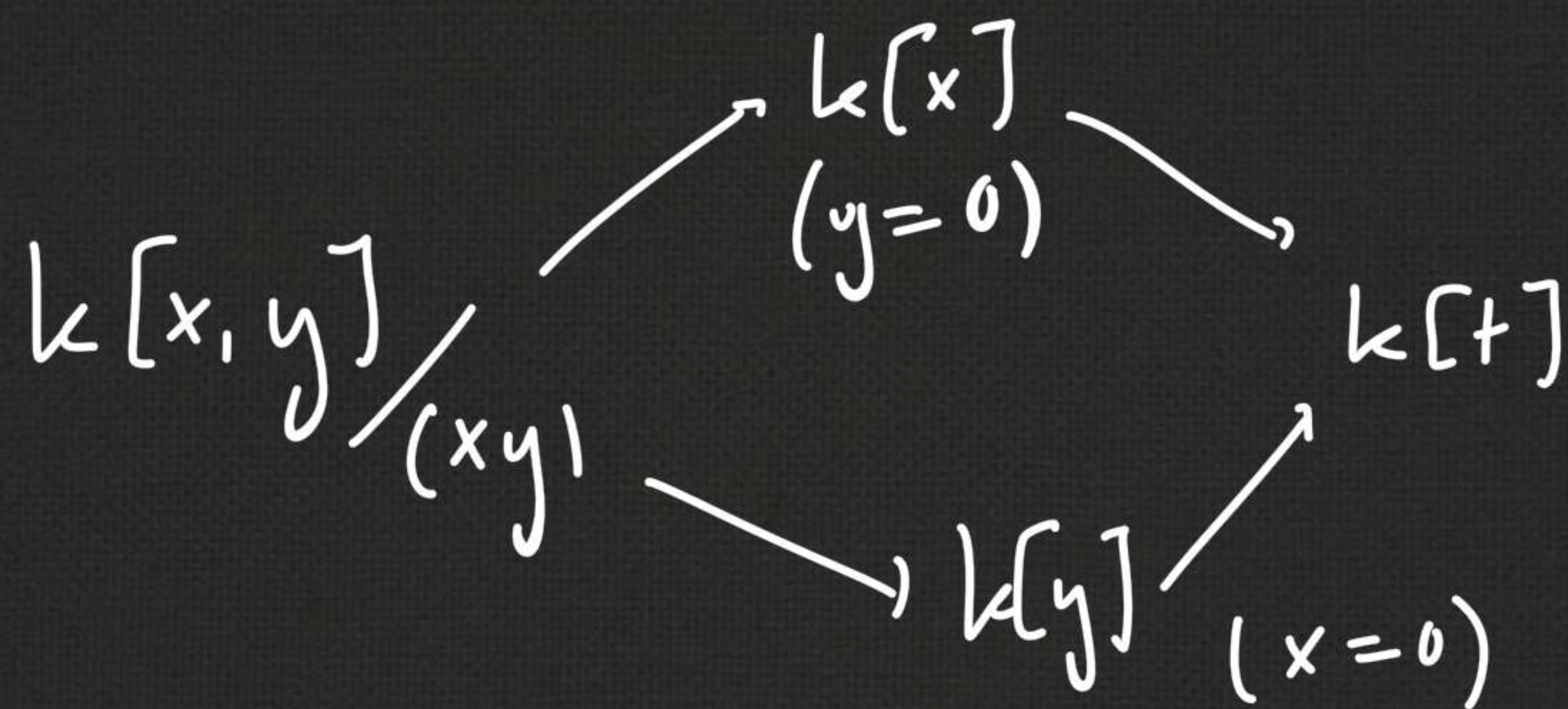
Idea: define two morphisms $X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$ to be " A' -homotopic" if $\exists h: X \times A' \rightarrow Y$ making the diagram



This almost works. If \sim is the relation that you get from this, then \sim is reflexive and symmetric, but not transitive.



Here $f \sim g$, and $g \sim h$, but $f \not\sim h$ since any $k[x,y]/(xy) \rightarrow k[t]$ factors



Defⁿ: an elementary finite correspondence from X to Y is an irreducible closed subset of $X \times Y$ whose associated integral scheme is finite and surjective over X via the projection $X \times Y \rightarrow X$. (e.g. if $f: X \rightarrow Y$, then Γ_f)

More generally, a finite correspondence from X to Y is an element of the free abelian group generated by elementary finite correspondences $\text{Cor}(X, Y)$.

For X, Y, Z smooth there is a composition

$$\text{Cor}(X, Y) \times \text{Cor}(Y, Z) \longrightarrow \text{Cor}(X, Z)$$

$$(\alpha, \beta) \longmapsto \beta \circ \alpha$$

defined on elementary finite correspondences $\alpha = [v], \beta = [w]$ as the pushforward cycle (along $X \times Y \times Z \rightarrow X \times Z$) of the intersection cycle $[v \times Z] \cap [X \times w]$.

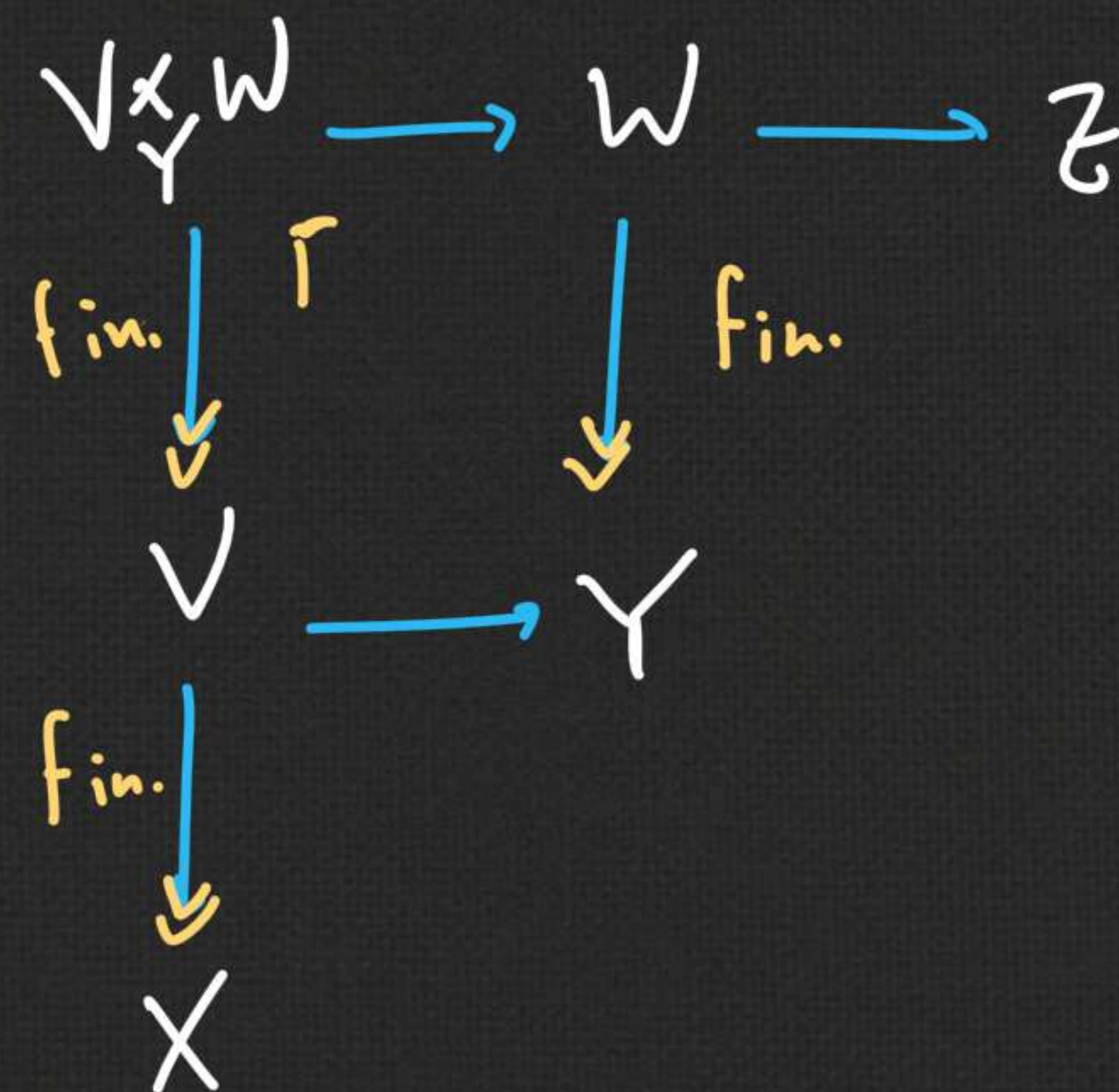
(if $f: X \rightarrow Y$ and $g: Y \rightarrow Z$, the $[\Gamma_g] \circ [\Gamma_f] = [\Gamma_{g \circ f}]$)

$$V \subset X \times Y \quad \text{and} \quad W \subset Y \times Z$$

$$V \times_Y W \rightarrow (X \times Y) \times_Y (Y \times Z) = X \times Y \times Z$$

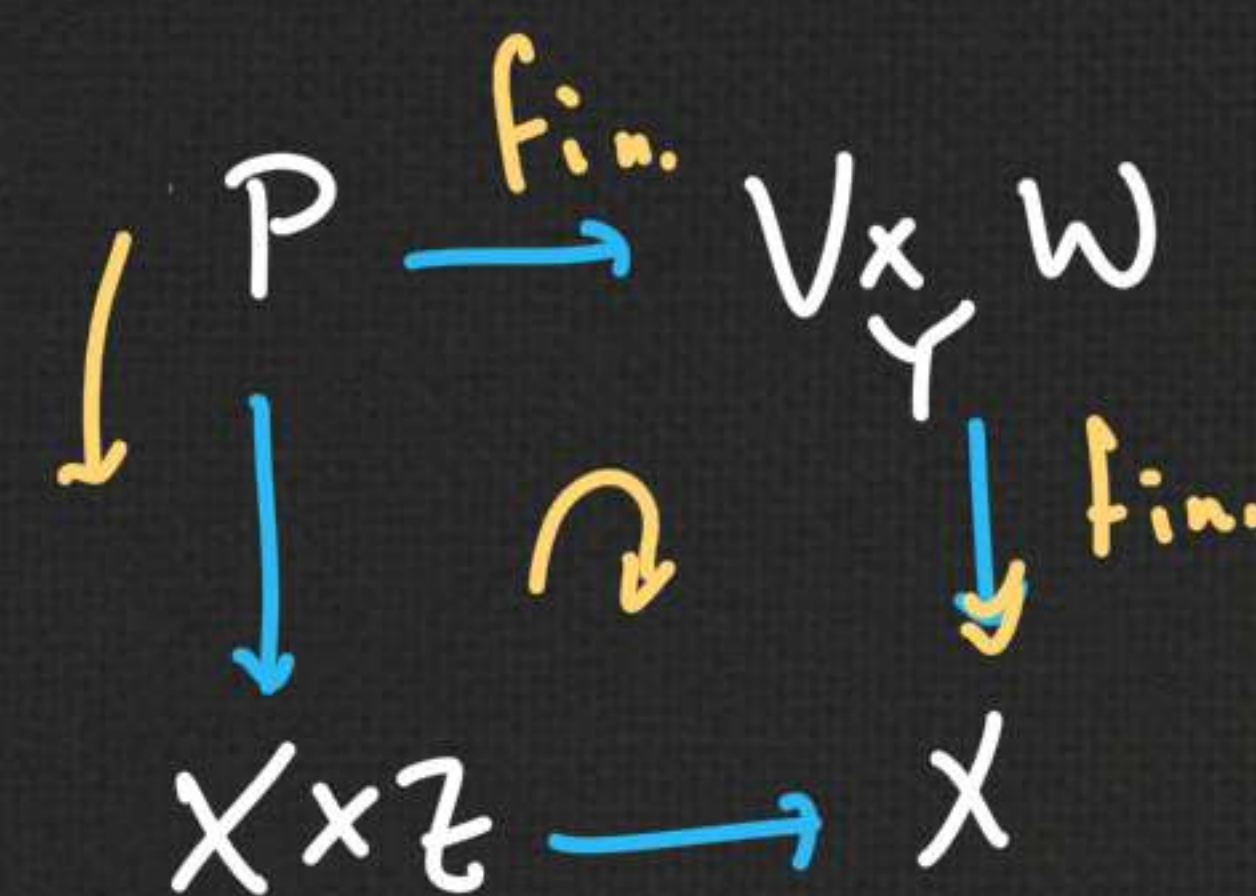
has image $(V \times Z) \cap (X \times W)$ and

$$\begin{aligned} \text{codim}(V \times_Y W) &= \dim(Z) + \dim(Y) \\ &= \text{codim}(X \times W) + \text{codim}(V \times Z). \end{aligned}$$



Then $[V \times Z] \cap [X \times W] = \sum_i m_p [P]$ where P irred.

$$m_p = \sum_{i \geq 0} (-1)^i \text{length}_{\mathcal{O}_{X \times Y \times Z, P}} T_{\text{or}_i}^{\mathcal{O}_{X \times Y \times Z, P}} \left(\mathcal{O}_{V \times Z, P}, \mathcal{O}_{X \times W, P} \right).$$



So $P \rightarrow X \times Z$ is finite, and $P \rightarrow X$.

Defⁿ: a naive A' -homotopy between morphisms $X \begin{array}{c} \xrightarrow{f^\alpha} \\ \xrightarrow{g^\beta} \end{array} Y$ is ^{or correspondences!}
 a finite correspondence $h \in \text{Cor}(X \times A', Y)$ so that

$$\left[[\Gamma_f^\alpha] = h \circ [\Gamma_{i_0}'] \quad \text{and} \quad [\Gamma_g^\beta] = h \circ [\Gamma_{i_1}'] \right]$$

where $X \begin{array}{c} \xrightarrow{i_0} \\ \xrightarrow{i_1} \end{array} X \times A'$ are the inclusions.

This is an equivalence relation (on $\text{Hom}(X, Y)$ or $\text{Cor}(X, Y)$)

that we'll denote by $[\sim_{A'}]$

Now with a working notion of (naive) A' -homotopies, we can define a notion of naive A' -homotopy equivalences.

Def^h: a morphism $f: X \rightarrow Y$ is a naive A' -homotopy equivalence if $\exists g: Y \rightarrow X$ so that $\Gamma_{g \circ f} \sim_{A'} \Delta_X$ and $\Gamma_{f \circ g} \sim_{A'} \Delta_Y$.

$\Gamma_{g \circ f} \stackrel{=}{=} \Gamma_{\text{id}_X} \subset X \times X$

Question: What properties of a variety are preserved by naive A' -homotopy equivalences?

Part 2:

a "rigidity" type theorem

Theorem 1: if X, Y are smooth projective varieties, then X, Y are naively \mathbb{A}^1 -homotopy equivalent if and only if $X \cong Y$.

In general, naive \mathbb{A}^1 -homotopy equivalences are nontrivial.

Ex For any smooth X , the projection $X \times \mathbb{A}^1 \rightarrow X$ is a naive \mathbb{A}^1 -homotopy equivalence.

Outline of Proof: (In a number of steps)

Step # 1: Show that two finite correspondences $\alpha, \beta \in \text{Cor}(X, Y)$ (with X, Y proper) define the same cycle class in $\text{CH}(X \times Y)$ if $\alpha \sim_{A'} \beta$.

Main difficulty:

$$\begin{array}{ccc}
 \text{Cor}(X \times A', Y) & \begin{array}{c} \xrightarrow{\circ \Gamma_{i_0}} \\ \xrightarrow{\circ \Gamma_{i_1}} \end{array} & \text{Cor}(X, Y) \\
 \downarrow & \circlearrowleft & \downarrow \\
 \text{CH}(X \times A' \times Y) & \begin{array}{c} \xrightarrow{\circ \Gamma_{i_0}} \\ \xrightarrow{\circ \Gamma_{i_1}} \end{array} & \text{CH}(X \times Y)
 \end{array}$$

Two Solutions

#1: lift from A' to P'

#2: show this diagram exists with Gysin pullbacks

Step #2: any cycle class $\alpha \in CH(X \times Y)$ defines a pushforward

$$\alpha_* : CH(X) \longrightarrow CH(Y)$$

$$\beta \longmapsto \pi_{Y*} (\alpha \wedge \pi_X^{-1}(\beta))$$

$$\begin{array}{ccc} & X \times Y & \\ \pi_X \swarrow & & \searrow \pi_Y \\ X & & Y \end{array}$$

So combined with Step #1, we get that if $\alpha \sim_{\mathbb{A}^1} \beta$

$$\text{then } \alpha_* = \beta_* : CH(X) \longrightarrow CH(Y).$$

(For naive \mathbb{A}^1 -homotopy equivalences $\Delta_X \sim_{\mathbb{A}^1} \Gamma_{g \circ f}$ and $\Delta_Y \sim_{\mathbb{A}^1} \Gamma_{f \circ g}$)

says $\Delta_X \cdot = \text{id}_{CH(X)} = [\Gamma_{g \circ f}] \cdot = (g \circ f)_* = \underline{g_* \circ f_*}$ and $\Delta_Y \cdot = \text{id}_{CH(Y)} = \underline{f_* \circ g_*}$.)

Ex Let $\mathbb{P}_{x,y}^1$ and $\mathbb{P}_{x,y,z}^2$. Then the map

$$h(t): \mathbb{P}^1 \times \mathbb{A}^1 \longrightarrow \mathbb{P}^2$$

$$([x:y], t) \longmapsto [x^2:txy:y^2]$$

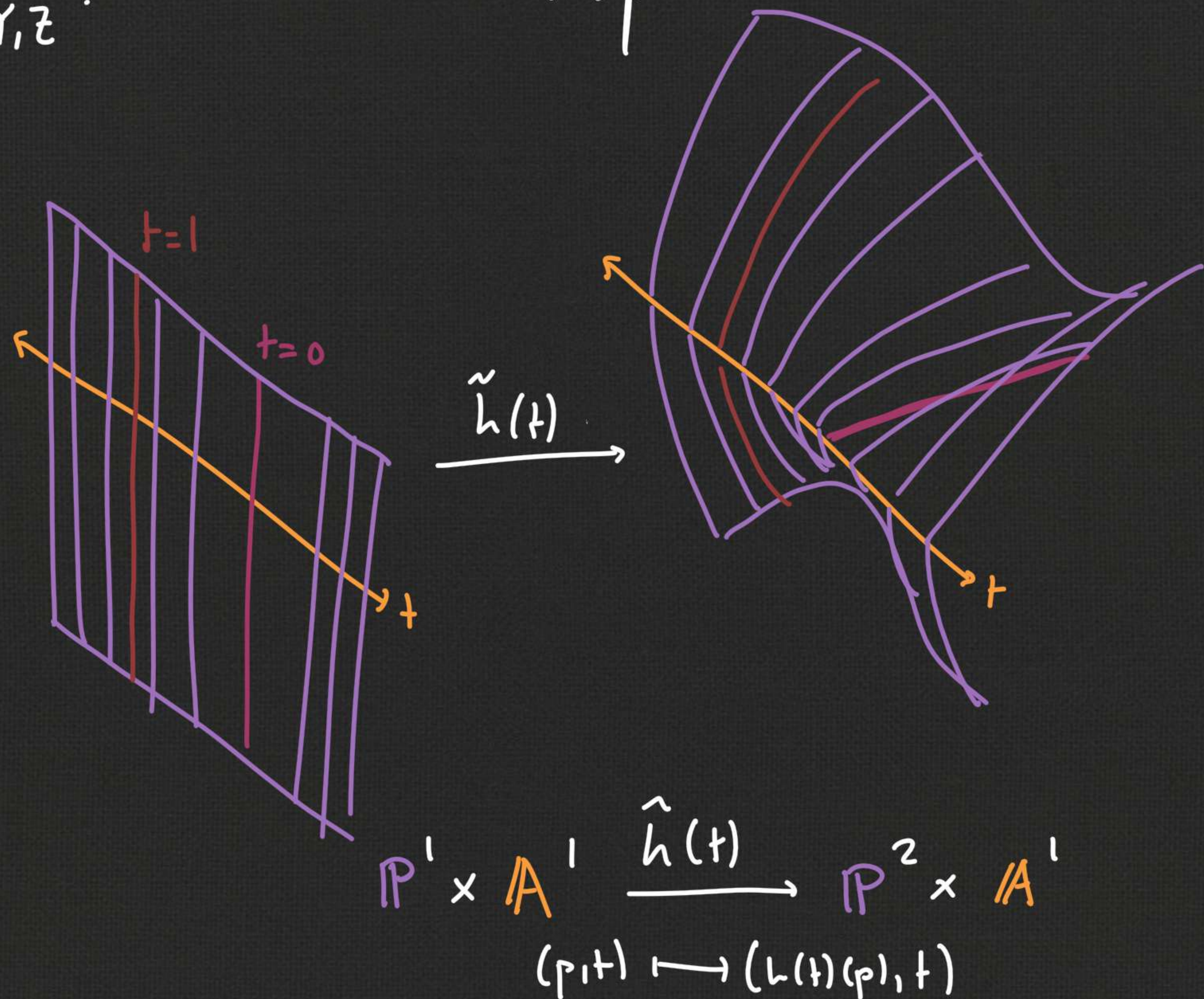
$$h(0)_*, h(1)_*: CH_0(\mathbb{P}^1) \xrightarrow{\sim} CH_0(\mathbb{P}^2)$$

$$h(0)_*: CH_1(\mathbb{P}^1) \longrightarrow CH_1(\mathbb{P}^2)$$

$$[\mathbb{P}^1] \longmapsto 2[\gamma=0]$$

$$h(1)_*: CH_1(\mathbb{P}^1) \longrightarrow CH_1(\mathbb{P}^2)$$

$$[\mathbb{P}^1] \longmapsto [xz=y^2]$$



Step #3: Prove a (strong) analog of a theorem of Whitehead:

Theorem 2: Let $f: X \rightarrow Y$ be a morphism between projective varieties X and Y such that

$$f_*: CH(X) \rightarrow CH(Y)$$

is an isomorphism. If Y is normal, then

f is an isomorphism.

(Altogether, this finishes the proof.)

✓

Part 3:

analogues of Whitehead's theorem

Whitehead's Theorem (together with Hurewicz's Theorem) says:

if $f: X \rightarrow Y$ is a continuous map between simply connected CW complexes with $f_*: H_*(X, \mathbb{Z}) \rightarrow H_*(Y, \mathbb{Z})$ an isomorphism, then f is a homotopy equivalence.

We prove a few variants of the previous Theorem also (in each case by reducing to Zariski's Main Theorem).

Theorem^H: if $f: X \rightarrow Y$ is a morphism (over $k = \mathbb{C}$) with X, Y projective smooth and $f_*: H_{2*}(X(\mathbb{C}), \mathbb{Z}) \rightarrow H_{2*}(Y(\mathbb{C}), \mathbb{Z})$ is an iso. Then so is f .

Theorem^K: if $f: X \rightarrow Y$ is a morphism with X, Y projective, Y normal, and $f_*: G(X) \rightarrow G(Y)$ an iso. Then so is f .

Theorem^Ω: if $f: X \rightarrow Y$ is a morphism with X, Y projective, Y normal, $\text{char}(k) = 0$, and if $f_*: \Omega_0(X) \rightarrow \Omega_0(Y)$ is an iso. Then so is f .

These theorems are sharp.

Ex (1) (Can't drop normality) If $C = \text{Proj}(k[x, y, z]/(y^2z = x^3))$ then the normalization $v: \mathbb{P}^1 \rightarrow C$ induces an isomorphism

$$\begin{array}{ccc} v_*: \text{CH}(\mathbb{P}^1) & \longrightarrow & \text{CH}(C) \\ \parallel & & \parallel \\ \mathbb{Z} \oplus \mathbb{Z} & & \mathbb{Z} \oplus \mathbb{Z} \end{array}$$

(2) (Can't use \mathbb{Q} coeff.) If Q is a 3 dim'l quadric and $k = \bar{k}$ then there is a finite map $Q \rightarrow \mathbb{P}^3$ with

$$\mathbb{Q}^{\oplus 4} = \text{CH}(Q)_{\mathbb{Q}} \longrightarrow \text{CH}(\mathbb{P}^3)_{\mathbb{Q}} = \mathbb{Q}^{\oplus 4} \quad \text{an isomorphism.}$$

Lemma: Let $f: X \rightarrow Y$ be a morphism with X, Y projective and assume $f_*: CH(X) \rightarrow CH(Y)$ has torsion kernel. Then f is finite.

Proof. Assume f is not finite. Let $\underline{V \subset X}$ have $\dim(V) > \dim(f(V)) \geq 0$. Then for any embedding $i: X \subset \mathbb{P}^n$ we get

$$\begin{array}{ccc}
 CH(X) & \xrightarrow{i_*} & CH(\mathbb{P}^n) \\
 f_* \downarrow & & \downarrow \\
 CH(Y) & & 0
 \end{array}
 \quad
 \begin{array}{ccc}
 [V] & \mapsto & \deg(V) [\mathbb{P}^{\dim(V)}] \neq 0 \\
 \downarrow & & \downarrow \\
 & & 0
 \end{array}$$

\nexists

Lemma: Let $f: X \rightarrow Y$ be a morphism with X and Y projective.

Assume that $f_*: CH(X) \rightarrow CH(Y)$ is surjective. Then

(1) f is surjective

(2) if f is finite, then f is birational.

Proof. For (1), let $j: U \subset Y$ be the inclusion of an open subscheme. Then

$$\begin{array}{ccc}
 f^{-1}(U) & \xrightarrow{\hat{f}} & X \\
 \hat{f} \downarrow & \uparrow & \downarrow f \\
 U & \xrightarrow{j} & Y
 \end{array}
 \rightsquigarrow
 \begin{array}{ccc}
 CH(f^{-1}(U)) & \xleftarrow{\hat{f}_*} & CH(X) \\
 \hat{f}_* \downarrow & & \downarrow f_* \\
 CH(U) & \xleftarrow{j_*} & CH(Y)
 \end{array}$$

Taking $U = Y \setminus f(X)$ shows $U = \emptyset$.

For (2), note

$$\begin{aligned}
 f_*: \mathbb{Z} = CH^0(X) &\longrightarrow CH^0(Y) = \mathbb{Z} \\
 [X]_1 &\longmapsto \deg(f) \cdot [f(X)].
 \end{aligned}$$

Altogether, if $f_*: CH(X) \rightarrow CH(Y)$ is an isomorphism, then

f is both finite and birational. If Y is normal, then

Zariski's Main Theorem implies f is an isomorphism. \square

THANKS!

